

Logics without the contraction rule and residuated lattices

HIROAKIRA ONO

JAPAN ADVANCED INSTITUTE OF SCIENCE AND TECHNOLOGY

NOMI, ISHIKAWA, JAPAN

ono@jaist.ac.jp

Received by Greg Restall

Published September 22, 2010

<http://www.philosophy.unimelb.edu.au/ajl/2010>

© 2010 Hiroakira Ono

To the memory of R. K. Meyer

Abstract: In this paper, we will develop an algebraic study of substructural propositional logics over FL_{ew} , i.e. the logic which is obtained from the intuitionistic logics by eliminating the contraction rule. Our main technical tool is to use residuated lattices as the algebraic semantics for them. This enables us to study different kinds of nonclassical logics, including intermediate logics, BCK-logics, Łukasiewicz's many-valued logics and fuzzy logics, within a uniform framework.

FOREWORD (2009): The draft of the present paper was originally completed in 1999 and then revised slightly in 2001, which I intended to dedicate to R. K. Meyer on the occasion of his 65th birthday.

From the middle of 90s, I had been trying to develop an algebraic study of substructural logics over the logic without the contraction rule FL_{ew} , and had announced results in several conferences, e.g. the 6th Asian Logic Conference in 1996, Dagstuhl Seminar on Multiple-Valued Logic in 1997 and S. Jaśkowski Memorial Symposium on Parainconsistent Logic, Logical Philosophy, Mathematics & Informatics in 1998. So I planned to make the present paper a comprehensive survey of the state of affairs of the study.

While the paper has not been published for many years, the draft has been referred in considerably many papers of substructural logics published in the last decade.

On the occasion of the publication of the paper, I have been wondering in which way I should publish the paper and how much I should revise it. For, there have been a remarkable progress in the direction of this research within these 10 years, and moreover my joint book [25] on substructural logics was already published. It would be no use to make an entire update of it.

My decision is to minimize changes, keeping the original form, and to add only necessary information on recent progress. (I put the mark (†) to footnotes which are essentially added in the present revision.) I hope that the paper will be still of interest and also informative, especially in its references, beyond its historical meaning.

I INTRODUCTION

A lot of works have been done in recent years for substructural logics, logics lacking some or all of structural rules when they are formalized in sequent calculi (see e.g. [21]). They include various kinds of nonclassical logics like Lambek calculus for categorial grammar, linear logic, BCK-logic and relevant logics, which have been introduced by different motivations with their own interests and aims and have been studied separately. The study of substructural logics, if successful, will enable us to develop the study of various nonclassical logics in a uniform viewpoint, and to discuss common features among them within this framework.

Though proof-theoretic methods have showed their effectiveness for particular substructural logics, e.g. logics formalized in cut-free sequent systems (see e.g. [54]), it will be quite necessary to introduce semantical methods when we want to develop a general study of substructural logics to the similar extent to that of modal logics.

In this paper, we will develop an algebraic study of substructural logics over FL_{ew} . The sequent calculus FL_{ew} is obtained from the intuitionistic logic by deleting the contraction rule. Sometimes we call *substructural logics over FL_{ew}* , or *logics without contraction rule*, although the contraction rule holds in some of them. The class of logics without the contraction rule contains intermediate logics, BCK-logics, Łukasiewicz's many-valued logics and fuzzy logics (in the sense of [30]). Our main technical tool here is algebraic one which is based on closer connections between logics over FL_{ew} and classes of residuated lattices.

Our main aims of the present paper are first to view known results on various logics without contraction in our framework, to keep the situation in perspective and to try to find out proper directions of further study of substructural logics over FL_{ew} . We give here some basic references of various classes of logics without contraction discussed in the present paper: [13] for intermediate logics (or superintuitionistic logics), [16] for many-valued logics, [30] for fuzzy logics, [25] for residuated lattices, [8, 9] for BCK-algebras and [21, 25] for substructural logics.

The paper is organized as follows. In Section 2, the logic FL_{ew} and its extensions are introduced. Results on residuated lattices and FL_{ew} -algebras are surveyed in Section 3. Basic results on subdirectly irreducible FL_{ew} -algebras are shown in Section 4. Some of important extensions of FL_{ew} are introduced in Section 5 as axiomatic systems. In Section 6, simple and semisimple FL_{ew} -algebras are discussed. A classification of logics over FL_{ew} is also introduced

at the end of the section. Section 7 is devoted to the study of immediate predecessors of the classical logic Cl. It is known that there is a single immediate predecessor of Cl over the intuitionistic logic, and that there are countably many of them over Łukasiewicz's infinitely many-valued logic. It will be shown there that there are many immediate predecessors of Cl other than them among logics over FL_{ew}.

The author owed much to Bob Meyer and Josep M. Font in completing the initial draft of the paper. The author would like to express his special thanks to them for constant encouragement and suggestions. The author is also indebted much to A. Avron, G. Bezhanishvili, W. Blok, M. Bunder, H. Ida, T. Kowalski, K. Swirydowicz and F. Wolter for their advices and helps in various forms.

2 LOGICS WITHOUT THE CONTRACTION RULE

We will introduce first a sequent calculus FL_{ew}, which is our basic logic without the contraction rule. Roughly speaking, FL_{ew} is the system obtained from Gentzen's sequent calculus LJ for the intuitionistic logic by deleting the contraction rule. The language of FL_{ew} consists of a logical constant \perp , logical connectives $\rightarrow, \wedge, \vee$ and \cdot (called *multiplicative conjunction* or *fusion*). The negation $\neg A$ of a formula A is defined as an abbreviation of $A \rightarrow \perp$. Sometimes, we will abbreviate the formula $(A \rightarrow B) \wedge (B \rightarrow A)$ to $A \equiv B$. A *sequent* is of the form $A_1, \dots, A_m \Rightarrow B$ where $m \geq 0$. In the following, capital Greek letters denote finite (possibly empty) sequences of formulas. The system FL_{ew} consists of the following initial sequents

1. $A \Rightarrow A$
2. $\perp, \Gamma \Rightarrow C$

and the following rules of inference;

Cut rule:

$$\frac{\Gamma \Rightarrow A \quad A, \Delta \Rightarrow C}{\Gamma, \Delta \Rightarrow C}$$

Exchange rule and weakening rule:

$$\frac{\Gamma, A, B, \Delta \Rightarrow C}{\Gamma, B, A, \Delta \Rightarrow C} \text{ (ex)} \qquad \frac{\Gamma \Rightarrow C}{A, \Gamma \Rightarrow C} \text{ (weak)}$$

Rules for logical connectives:

$$\frac{A, \Gamma \Rightarrow B}{\Gamma \Rightarrow A \rightarrow B} \text{ } (\Rightarrow \rightarrow) \qquad \frac{\Gamma \Rightarrow A \quad B, \Delta \Rightarrow C}{A \rightarrow B, \Gamma, \Delta \Rightarrow C} \text{ } (\rightarrow \Rightarrow)$$

$$\frac{\Gamma \Rightarrow A}{\Gamma \Rightarrow A \vee B} \text{ } (\Rightarrow \vee 1) \qquad \frac{\Gamma \Rightarrow B}{\Gamma \Rightarrow A \vee B} \text{ } (\Rightarrow \vee 2)$$

$$\frac{A, \Gamma \Rightarrow C \quad B, \Gamma \Rightarrow C}{A \vee B, \Gamma \Rightarrow C} (\vee \Rightarrow)$$

$$\frac{\Gamma \Rightarrow A \quad \Gamma \Rightarrow B}{\Gamma \Rightarrow A \wedge B} (\Rightarrow \wedge)$$

$$\frac{A, \Gamma \Rightarrow C}{A \wedge B, \Gamma \Rightarrow C} (\wedge 1 \Rightarrow) \quad \frac{B, \Gamma \Rightarrow C}{A \wedge B, \Gamma \Rightarrow C} (\wedge 2 \Rightarrow)$$

$$\frac{\Gamma \Rightarrow A \quad \Delta \Rightarrow B}{\Gamma, \Delta \Rightarrow A \cdot B} (\Rightarrow \cdot) \quad \frac{A, B, \Gamma \Rightarrow C}{A \cdot B, \Gamma \Rightarrow C} (\cdot \Rightarrow).$$

The provability of a given sequent is defined in the usual way. In particular, we say that a formula A is provable in FL_{ew} when the sequent $\Rightarrow A$ is provable in it.

In our joint paper [55] with Y. Komori, both syntactic and semantic properties of FL_{ew} are studied. The cut elimination theorem for FL_{ew} is shown, from which both the decidability and Craig's interpolation theorem of FL_{ew} are derived. Also, a Hilbert-style formulation of FL_{ew} is introduced and the separation theorem is proved. Then, a Kripke-type semantics for FL_{ew} and related systems is introduced and their completeness with respect to the semantics is proved in it. (For additional information on FL_{ew} and related systems (up to the middle of 90s), see also [22, 23, 52, 53].) It should be also noticed that *monoidal propositional logic* in [33] discussed in the context of fuzzy logic is equivalent to the logic FL_{ew} .

Let FL_e be the sequent system obtained from FL_{ew} by deleting the weakening rule, and FL_{ec} be the sequent system obtained from FL_e by adding the following contraction rule:

$$\frac{A, A, \Gamma \Rightarrow C}{A, \Gamma \Rightarrow C} (\text{con}).$$

The system FL_e gives a sequent calculus for the intuitionistic linear logic. It is clear that the sequent calculus obtained from FL_{ew} by adding the contraction rule gives a (cut-free) sequent calculus for the intuitionistic logic Int . We can see that any sequent of the form $A \cdot B \Rightarrow A \wedge B$ (and of the form $A \wedge B \Rightarrow A \cdot B$) is provable in FL_{ew} (and FL_{ec} , respectively). Thus in Int , the fusion \cdot becomes equivalent to \wedge .

In the present paper, we will concentrate on the study of substructural propositional logics over FL_{ew} . Here, by a *substructural logic over FL_{ew}* (or a *substructural logic without the contraction rule*), we mean any set of formulas which includes all formulas provable in FL_{ew} and is closed under substitution and *modus ponens*, i.e. if both A and $A \rightarrow B$ belong to the set then B belongs also to it. Sometimes, we omit the word "substructural", or even call it simply a *logic*.¹ Also, we identify a formal system with the set of all formulas which are

¹(†) As for the definition of *substructural logic* in general, see [25].

provable in it, when no confusions will occur. Thus, FL_{ew} denotes not only the sequent calculus introduced in the above, but also the set of all formulas provable in it. In the following, Cl and Int denote classical logic and intuitionistic logic, respectively. Intermediate logics (or, superintuitionistic logics), which are logics over intuitionistic logic, Łukasiewicz's many-valued logics, and fuzzy logics are important examples of logics over FL_{ew} , which have been already studied extensively.

The class of logics over FL_{ew} is ordered by the set inclusion \subseteq . (It is in fact a set, since each logic is a subset of the set of all formulas.) Of course, FL_{ew} is the smallest logic among them, and the set of all formulas is the greatest one, which is the *inconsistent* logic. We are concerned only with *consistent* logics, among which the classical logic Cl is the greatest.

Suppose that $\{L_i\}_{i \in I}$ is a set of logics, where I is a (possibly infinite) nonempty set of indices. Then, clearly the set intersection $\bigcap_{i \in I} L_i$ of them is also a logic. Thus, the set of all logics over FL_{ew} forms a complete lattice, in which the join $\bigvee_{i \in I} L_i$ of L_i s is represented as follows.

$$\bigvee_{i \in I} L_i = \{A : \text{there exist } j_1, \dots, j_k \in I \text{ and formulas } B_{j_t} \in L_{j_t} \text{ for } 1 \leq t \leq k \text{ such that the formula } (B_{j_1} \cdot \dots \cdot B_{j_k}) \rightarrow A \text{ is provable in } FL_{ew}\}.$$

THEOREM 2.1 *The set of all logics over FL_{ew} forms a complete lattice, in which the following distributive law hold;*

$$L \cap \bigvee_i L_i = \bigvee_i (L \cap L_i).$$

The distributivity in the above theorem can be easily shown by using the above representation of joins. We can derive it also from the fact that *the variety of FL_{ew} -algebras is congruence-distributive* (see Proposition 3.5), by using the result by Jónsson [39] (see also [5]).

Let L_0 and L be logics such that $L_0 \subseteq L$. Then, L is said to be *finitely axiomatized over L_0 by the axioms A_1, \dots, A_m* , if L is the smallest logic which contains both L_0 and the set $\{A_1, \dots, A_m\}$. It holds that for any formula C it is in L if and only if there exist formulas B_1, \dots, B_n (for some $n \geq 0$), each of which is a substitution instance of some A_k , such that the formula $(B_1 \cdot \dots \cdot B_n) \rightarrow C$ belongs to L_0 . The logic L is denoted by $L_0[A_1, \dots, A_m]$ in this case. A logic L is said to be *finitely axiomatizable over L_0* when there exist some axioms by which L is finitely axiomatized over L_0 . We will omit the word "over L_0 " when L_0 is FL_{ew} . It is easy to see that $L[A_1, \dots, A_m] = L[(A_1 \cdot \dots \cdot A_m)]$, by the help of the *weakening rule* of FL_{ew} , i.e. by using the fact that formulas of the form $(C \cdot D) \rightarrow C$ is always provable in FL_{ew} . It is easy to see the following.

THEOREM 2.2 *Suppose that logics L and L' are finitely axiomatized over L_0 by the axioms A and B , respectively. Then, logics $L \wedge L'$ and $L \vee L'$ are finitely axiomatized over L_0 by axioms $A \vee B$ and $A \cdot B$, respectively. (Here, A and B are supposed to have no propositional variables in common, by renaming them if necessary.)*

3 RESIDUATED LATTICES AND FL_{ew} -ALGEBRAS

In this section, we will introduce FL_{ew} -algebras as the algebraic counterparts of logics without contraction rules. FL_{ew} -algebras are sometimes called *commutative integral residuated lattices*. Residuated lattices are already introduced and discussed in 30s and has been studied by many people, e.g. Krull [47], Balbes-Dwinger [4], Dilworth [20], Ward-Dilworth [63], Ward [62] and Pavelka [58].²

In our original draft of the present paper in 2001, residuated lattices are defined in a slightly different way. In fact, they should have been called *commutative integral residuated lattices*, or more precisely FL_{ew} -algebras (see [53]). To avoid confusion, here we start to give a definition of *commutative residuated lattices* and then give a definition of FL_{ew} -algebras. As for a definition of residuated lattices in general, including noncommutative case, see [25].³

DEFINITION 3.1 An algebra $M = \langle M, \wedge, \vee, \cdot, \rightarrow, 1 \rangle$ is a commutative residuated lattice if

- $\langle M, \wedge, \vee \rangle$ is a lattice,
- $\langle M, \cdot, 1 \rangle$ is a commutative monoid with the unit element 1,
- for $x, y \in M$, $x \cdot y \leq z$ if and only if $x \leq y \rightarrow z$ (the law of residuation).

DEFINITION 3.2 An algebra $M = \langle M, \wedge, \vee, \cdot, \rightarrow, 0, 1 \rangle$ is a FL_{ew} -algebra if

- $\langle M, \wedge, \vee, \cdot, \rightarrow, 1 \rangle$ is a commutative residuated lattice,
- 1 is the greatest element and 0 is the least element.

FL_{ew} -algebras are called also *BCK-lattices* [34], *full BCK-algebras* [55], and *integral, residuated, commutative l-monoids* [33]. An extensive study of FL_{ew} -algebras can be seen in [33]. The (dual of the) *implicational reduct* of an FL_{ew} -algebra is sometimes called a *BCK-algebra*. A BCK-algebra with the fusion satisfying the law of residuation is called a BCK-algebra with condition (S), in [35]. See also [65], [9] and [10].

² In [67], Zlatoš discussed residuated lattices and claimed that the class of residuated lattices forms a variety which is arithmetical. Unfortunately, his definition of residuated lattices is insufficient. In fact, from his definition we cannot deduce the law of residuation in Definition 3.2 of residuated lattices. Thus, Proposition 1.2 of [67] is incorrect. The present author owes this fact to J. M. Font.

³ By the abuse of symbols, we use the same symbols both for logical connectives and for algebraic operations corresponding to these connectives.

Now we will give a brief survey of basic properties of FL_{ew} -algebras and of relations between logics over FL_{ew} and FL_{ew} -algebras.⁴ In the following, we assume that FL_{ew} -algebras are always *non-degenerate*, i.e. they satisfies $0 \neq 1$. We define $\sim x$ by $\sim x = x \rightarrow 0$. It is easy to see that the following hold in any FL_{ew} -algebra (in fact, the first three hold in any commutative residuated lattice);

- (i) $x \leq y$ implies $x \cdot z \leq y \cdot z$,
- (ii) $(x \vee y) \cdot z = (x \cdot z) \vee (y \cdot z)$,
- (iii) $1 \rightarrow x = x$,
- (iv) $x \cdot y \leq x$ and hence $x \cdot y \leq x \wedge y$ (integrality),
- (v) $x \leq y$ if and only if $x \rightarrow y = 1$.

Note that the above (ii) expresses the distributivity of \cdot over join, which follows from (i) with the law of residuation. Conversely, suppose that the distributivity (ii) holds in a *finite* algebra M satisfying Conditions 1 and 2 in the definition of commutative residuated lattices. Then, M becomes a commutative residuated lattice if we define $y \rightarrow z$ by $\max\{x : x \cdot y \leq z\}$. In fact, it exists always for all y, z , and \rightarrow satisfies the law of residuation.

In the previous section, we mentioned briefly a relationship between logics over FL_{ew} and fuzzy logic. We will discuss it here in more detail. In fuzzy set theory, the set of *truth values* is the unit interval $[0, 1]$, which is linearly ordered by the natural order. A binary function T from $[0, 1] \times [0, 1]$ to $[0, 1]$ is called a *triangular norm* (simply, a *t-norm*) in the theory of *probabilistic metric spaces* (see e.g. [60]), if the following holds for x, y, z ;

- 1. $T(x, T(y, z)) = T(T(x, y), z)$,
- 2. $T(x, y) = T(y, x)$,
- 3. $T(x, 1) = x$,
- 4. $x \leq y$ implies $T(x, z) \leq T(y, z)$.

Thus, $\langle [0, 1], \cdot, 1 \rangle$ forms a commutative monoid in which (i) holds, if we define \cdot by $x \cdot y = T(x, y)$. Now we suppose moreover that T is a *continuous t-norm*, i.e. a t-norm which is a continuous function over the interval $[0, 1]$. For any y, z , define a subset $I_{y,z}$ of $[0, 1]$ by $I_{y,z} = \{x : T(x, y) \leq z\}$ and let $u = \sup I_{y,z}$. By using the continuity of T , we have that $T(u, y) = \sup\{T(x, y) : x \in I_{y,z}\} \leq z$. Hence, u is in fact equal to $\max I_{y,z}$. Thus, the law of residuation holds also, if we define $y \rightarrow z$ by $\max I_{y,z}$. Therefore, the unit interval with any continuous

⁴(†) By a rapid development of the study in recent years, most of results in this section may be standard now.

t-norm forms an FL_{ew} -algebra. This gives us a reason why fuzzy logics (in the narrow sense) can be regarded as logics over FL_{ew} (see also [30]).⁵

Note that in the above only the *left-continuity* of T , i.e. $T(\bigvee_i x_i, y) = \bigvee_i T(x_i, y)$, is used for showing the existence of *residuation* \rightarrow . It is easily seen that when \cdot is defined by $x \cdot y = T(x, y)$ for all x, y with a left-continuous t-norm T , the structure $\langle [0, 1], \cdot, 1 \rangle$ is nothing but an *integral, unital commutative quantale* with the universe $[0, 1]$. (See [59] for more information on quantales.)⁶

As we will show later (see Proposition 3.3), the class of all FL_{ew} -algebras *determines* the logic FL_{ew} . An FL_{ew} -algebra M is said to be *involutive* if

$$DN: \sim\sim x = x \text{ for any } x$$

holds always in it. Grišin discussed properties of involutive FL_{ew} -algebras in [27, 28, 29], where they are called *lattice- L° -algebras* (see also [51], in which they are called *Grišin algebras*). Also, involutive FL_{ew} -algebras have been studied by E. Casari and P. Minari as algebras for a *comparative logic*, and they are called *Abelian lattice-ordered zeroids*. (See a survey of comparative logic in Casari [12].) We can show the following.

LEMMA 3.1 *In any involutive FL_{ew} -algebra, the following holds.*

1. $\sim(x \vee y) = \sim x \wedge \sim y$,
2. $x \cdot y = \sim(x \rightarrow \sim y)$.

FL_e -algebras are defined similarly to FL_{ew} -algebras, but by deleting the second condition of Definition 3.2, which says that 1 is the greatest element and 0 is the least element. (Thus, 0 is an arbitrary element.) It is well-known that the class of involutive FL_e -algebras determines the linear logic (without exponentials) MALL introduced by Girard [26]. Complete involutive FL_e -algebras (i.e. *complete* as lattices) are called *unital, commutative quantales* (see e.g. [59]).

The algebraic condition which corresponds to the contraction rule is $z \leq z \cdot z$ (square-increasingness). From this, the inequality $x \wedge y \leq x \cdot y$ follows. Thus, for each FL_{ew} -algebra M , the monoid operation \cdot of M is square-increasing, if and only if the monoid operation \cdot of M is equal to \wedge , if and only if M is a Heyting algebra. Another interesting class of algebras related to FL_{ew} -algebras arose from Łukasiewicz's many-valued logics. Algebras in the class are called by various names, e.g. *MV-algebras* in [14], *Wajsberg algebras* in [24], *CN-algebras* in [41] and (*bounded*) *commutative BCK-algebras* in [66, 36] etc.. (As for the exact relation between MV-algebras and Wajsberg algebras, see e.g. [24].) We will give here a short sketch of them and their connection with FL_{ew} -algebras.

⁵See [2] and [17] for important results on the logic of continuous t-norms up to 2001.

⁶(†) See also [25].

DEFINITION 3.3 An algebra $W = \langle W, \rightarrow, \sim, 1 \rangle$ is a Wajsberg algebra if

- $1 \rightarrow x = x$,
- $(x \rightarrow y) \rightarrow ((y \rightarrow z) \rightarrow (x \rightarrow z)) = 1$,
- $(x \rightarrow y) \rightarrow y = (y \rightarrow x) \rightarrow x$,
- $(\sim x \rightarrow \sim y) \rightarrow (y \rightarrow x) = 1$.

In a Wajsberg algebra W , let us define the binary relation \leq on W by $x \leq y$ if and only if $x \rightarrow y = 1$. Then $\langle W, \leq \rangle$ is a partially ordered set with the greatest element 1 , and the least element 0 if we define 0 by ~ 1 . Moreover, if we define $x \vee y = (x \rightarrow y) \rightarrow y$ and $x \wedge y = \sim(\sim x \vee \sim y)$ then $x \vee y$ and $x \wedge y$ are equal to the *supremum* and the *infimum* of $\{x, y\}$, respectively, with respect to the order \leq . Hence W forms a lattice. Also, it can be shown that $x \rightarrow 0 = \sim x$, $\sim \sim x = x$ and moreover that $(x \rightarrow y) \vee (y \rightarrow x) = 1$ (prelinearity). The following equation (i), which is the third condition in Definition 3.3, plays an important role as shown below.

$$(x \rightarrow y) \rightarrow y = (y \rightarrow x) \rightarrow x. \tag{i}$$

It is easy to see that any FL_{ew} -algebra satisfying (i) is involutive. (As for the details of the above discussions, see e.g. [24]. See also [8] for related topics.) The following proposition is proved in [24] and [49].

PROPOSITION 3.1 *Let $W = \langle W, \rightarrow, \sim, 1 \rangle$ be a Wajsberg algebra. Define a binary relation \leq on W as above, and a binary operation \cdot by $x \cdot y = \sim(x \rightarrow \sim y)$. Then, $W = \langle W, \wedge, \vee, \cdot, \rightarrow, 0, 1 \rangle$ forms a (involutive) FL_{ew} -algebra. Conversely, each FL_{ew} -algebra satisfying $(x \rightarrow y) \rightarrow y = (y \rightarrow x) \rightarrow x$ naturally determines a Wajsberg algebra.*

The above proposition assures us that Wajsberg algebras can be regarded as FL_{ew} -algebras. Wajsberg algebras are originally introduced as models of Łukasiewicz's many-valued logics. Let us define the following operations either on the set $\{0, 1/n, 2/n, \dots, (n-1)/n, 1\}$ for a positive integer n or on the unit interval $[0, 1]$.

- $x \rightarrow y = \min\{1, 1 - x + y\}$,
- $\sim x = x \rightarrow 0 = 1 - x$,
- $x \cdot y = \sim(x \rightarrow \sim y) = \max\{0, x + y - 1\}$.

Then in either case, it forms a Wajsberg algebra. They are denoted as L_{n+1} and L , respectively, in the present paper. We have also the following (see also Lemmas 2.14 and 2.5 of [33]).

PROPOSITION 3.2 *For any FL_{ew} -algebra M , M satisfies the equation (i) if and only if it is involutive and satisfies $x \wedge y = x \cdot (x \rightarrow y)$.*

Proof: The only-if part can be obtained by using equations in Lemma 3.1. The if-part holds, since $x \vee y = \sim(\sim x \wedge \sim y) = (\sim x \cdot (\sim x \rightarrow \sim y)) \rightarrow 0 = ((y \rightarrow x) \cdot \sim x) \rightarrow 0 = (y \rightarrow x) \rightarrow (\sim x \rightarrow 0) = (y \rightarrow x) \rightarrow x$. Similarly, $x \vee y = y \vee x = (x \rightarrow y) \rightarrow y$. Thus, equation (1) holds. \square

Any FL_{ew} -algebra satisfying the equation $(x \rightarrow y) \vee (y \rightarrow x) = 1$ is called an MLT-algebra. An MLT-algebra satisfying the equation $x \wedge y = x \cdot (x \rightarrow y)$ is called a BL-algebra (see [30]). By Proposition 3.2, Wajsberg algebras are (essentially) equal to involutive BL-algebras.

In the usual way, we will define the *validity* of formulas (of FL_{ew}) in a given FL_{ew} -algebra M as follows. Any mapping v from the set of all propositional variables to the set M is called a *valuation* on M . A given valuation v can be extended to a mapping from the set of all formulas to M , inductively as follows.

1. $v(\perp) = 0$,
2. $v(A \wedge B) = v(A) \wedge v(B)$,
3. $v(A \vee B) = v(A) \vee v(B)$,
4. $v(A \cdot B) = v(A) \cdot v(B)$,
5. $v(A \rightarrow B) = v(A) \rightarrow v(B)$.

A formula A is *valid* in M if $v(A) = 1$ holds for any valuation on M . The set of formulas which are valid in M is denoted by $L(M)$. Next, a given sequent $A_1, \dots, A_m \Rightarrow B$ is said to be valid in M if the formula $(A_1 \cdot \dots \cdot A_m) \rightarrow B$ is valid in it. Then, the following completeness theorem of FL_{ew} can be shown in a standard way.

PROPOSITION 3.3 *A sequent S is provable in FL_{ew} if and only if it is valid in all FL_{ew} -algebras.*

It is easy to see that $L(M)$ is a logic over FL_{ew} for any FL_{ew} -algebra M , which is called the logic *determined* by M . (In general, for any class \mathcal{K} of FL_{ew} -algebras, the logic $\bigcap_{M \in \mathcal{K}} L(M)$ is a logic over FL_{ew} , called the logic determined by \mathcal{K} .) Conversely, for any logic L over FL_{ew} there exists an FL_{ew} -algebra M such that $L = L(M)$. In fact, it is enough to take the *Lindenbaum algebra* of L for M .

Blok and Ferreirim have developed the study of algebras called *hoops*, in [6, 7]. Each hoop satisfies the equation $x \cdot (x \rightarrow y) = y \cdot (y \rightarrow x)$. When a hoop satisfies moreover the equation (1), it is called a *Wajsberg hoop*. Any FL_{ew} -algebra with the equation (1) is an example of a Wajsberg hoop. As for the details, see [6, 7].

In [64], Wroński proved that the class of all BCK-algebras does not form a variety.⁷ (See also Higgs [32].) On the other hand, Idziak [34] proved the following.

⁷(†) For basic notations and results on universal algebra, see [11].

PROPOSITION 3.4 *The class of all FL_{ew} -algebras forms a variety.*

A variety \mathcal{K} is *congruence-distributive* when the lattice of all congruence relations of any algebra in \mathcal{K} is distributive, and is *congruence permutable* when every two congruences of any algebra in \mathcal{K} permute. Moreover, if \mathcal{K} is both congruence-distributive and congruence permutable, it is said to be *arithmetical*. Then the following holds (see [34]).

PROPOSITION 3.5 *The variety of FL_{ew} -algebras is arithmetical.*

Next, we will give the definition of filters of FL_{ew} -algebras.⁸

DEFINITION 3.4 A nonempty subset F of an FL_{ew} -algebra M is an *implicative filter* (or, simply a *filter*) if for $x, y \in M$ it satisfies

- $x \leq y$ and $x \in F$ imply $y \in F$,
- $x, y \in F$ implies $x \cdot y \in F$.

Note that any filter F of an FL_{ew} -algebra satisfies that $x, y \in F$ implies $x \wedge y \in F$. It is easy to see that a nonempty subset F of a residuated lattice is a filter if and only if it satisfies that 1) $1 \in F$ and 2) $x, x \rightarrow y \in F$ implies $y \in F$ (see e.g. [24]). Let S be a nonempty subset of an FL_{ew} -algebra M . Then the set $\{x : a_1 \cdots a_k \leq x \text{ for some } a_1, \dots, a_k \in S\}$ is shown to be a filter, called the filter *generated* by S . In particular, the filter generated by a singleton set $\{a\}$ for an element $a \in M$ will be expressed as $\{x : a^k \leq x \text{ for some positive integer } k\}$, where a^k denotes $a \cdots a$ with k times a . The following can be shown easily (see e.g. [33]).

PROPOSITION 3.6 *Let M be any FL_{ew} -algebra. Then, there exists a lattice isomorphism between the set of all filters of M and the set of all congruences of M .*

In fact, for a given filter F the binary relation \sim_F defined by

$$x \sim_F y \text{ if and only if } x \rightarrow y, y \rightarrow x \in F$$

is a congruence and the map θ defined by $\theta(F) = \sim_F$ is the required map. The converse map is defined by taking the filter $\{x : x \sim 1\}$ for a given congruence \sim .

4 SUBDIRECTLY IRREDUCIBLE FL_{ew} -ALGEBRAS

Let M and N_i for each $i \in I$ be residuated lattices. By a *subdirect representation* of M with factors N_i we mean an embedding f from M to the product $\prod_{i \in I} N_i$ such that each f_i defined by $f_i = p_i \circ f$ is onto N_i for each $i \in I$. Here, p_i denotes the i -th projection. An FL_{ew} -algebra M is *subdirectly irreducible* if it is non-degenerate and for any subdirect representation $f : M \rightarrow \prod_{i \in I} N_i$, there exists a j such that f_j is an isomorphism of M onto N_j (see [48] for the details).

⁸(†) In [25], they are called *deductive filters*. For the definition of filters in arbitrary residuated lattices, see [25]. Filter generation and its relation to (generalized form of) deduction theorem are also discussed in it.

From Birkhoff's *subdirect representation theorem* it follows that every FL_{ew} -algebra has a subdirect representation with subdirectly irreducible FL_{ew} -algebras. By using Proposition 3.6, we can see that an FL_{ew} -algebra M is subdirectly irreducible if and only if it has the second smallest filter, i.e. the smallest filter among all filters except $\{1\}$. Note that when an FL_{ew} -algebra M is subdirectly represented by the set $\{N_j\}_{j \in J}$ of FL_{ew} -algebras, the logic $L(M)$ determined by M can be expressed as $\bigcap_{j \in J} L(N_j)$.

We can easily show the following lemma which gives a necessary and sufficient condition for an FL_{ew} -algebra to be subdirectly irreducible.

LEMMA 4.1 *An FL_{ew} -algebra M is subdirectly irreducible if and only if there exists an element $a (< 1)$ such that for any $x < 1$ there exists a positive integer m for which $x^m \leq a$ holds.*

Using this lemma, we can show the following.

THEOREM 4.1 *In any subdirectly irreducible FL_{ew} -algebra, if $x \vee y = 1$ then either $x = 1$ or $y = 1$ holds.*

Proof: By taking the contraposition, it suffices to show that $x, y < 1$ implies $x \vee y < 1$ in a given subdirectly irreducible FL_{ew} -algebra M . Since M is subdirectly irreducible, there exists $a < 1$ such that for any $z < 1$ there exists a number k satisfying $z^k \leq a$. In particular, both $x^m \leq a$ and $y^n \leq a$ hold for some positive integers m and n . Define $s = \max\{m, n\}$ and $t = 2s - 1$. Then, clearly $x^s \leq a$ and $y^s \leq a$ hold. Now, by the distributivity of \cdot with \vee ,

$$(x \vee y)^t = \bigvee_{i=1}^t x^i \cdot y^{t-i}.$$

It is easy to see that either $i \geq s$ or $t - i \geq s$. Hence, in the former case,

$$x^i \cdot y^{t-i} \leq x^i \leq x^s \leq a$$

and in the latter case,

$$x^i \cdot y^{t-i} \leq y^{t-i} \leq y^s \leq a.$$

Thus, in either case, $(x \vee y)^t \leq a$. Therefore, $x \vee y$ cannot be equal to 1.

An element a in an FL_{ew} -algebra M is a *coatom* if it is maximal among elements in $M - \{1\}$. Then we have the following immediately from the above theorem.

COROLLARY 4.1 *Every subdirectly irreducible FL_{ew} -algebra has either the single coatom or no coatoms.*

The following result is essentially due to Kowalski [42], which makes an interesting contrast with Lemma 4.1.

LEMMA 4.2 *An FL_{ew} -algebra M has the unique coatom if and only if there exists an element $a (< 1)$ and a positive integer m such that $x^m \leq a$ holds for any $x < 1$.*

Proof: To show the only-if part, it is enough to take the coatom for a and 1 for m . Conversely, suppose that there exists $a (< 1)$ and $m \geq 1$ such that $x^m \leq a$ for any $x < 1$. By Lemma 4.1, it is clear that M is subdirectly irreducible. Now, take any such a and take also the smallest number k among such m s for this a . If $k = 1$ then it is obvious that a is the single coatom of M . So, suppose that $k > 1$. By our assumption, there exists an element b such that $b^{k-1} \not\leq a$ but $b^k \leq a$. Define d by $d = b^{k-1} \rightarrow a$. Clearly, $d < 1$. We will show that d is the single coatom of M . Take any y such that $y < 1$ and let $z = d \vee y$. Then, $d, y \leq z$ and moreover $z < 1$ by Theorem 4.1. Since $b^k \leq a$, $b \leq b^{k-1} \rightarrow a = d \leq z$. As $z < 1$, we have $z^k \leq a$ by our assumption. Thus, $y \leq z \leq z^{k-1} \rightarrow a \leq b^{k-1} \rightarrow a = d$. Therefore, $y \leq d$. Hence, d is the coatom of M . \square

Let us consider the following condition on FL_{ew} -algebras for a given positive integer k :

$$E_k : x^{k+1} = x^k \text{ for any } x.$$

The condition E_k is equivalent to the condition that $x^{k+1} \rightarrow y = x^k \rightarrow y$ for all x, y (see e.g. [7]). It is introduced and discussed by Cornish [18]. Sometimes, M is said to be *k-potent* when it satisfies E_k . It is obvious that every finite FL_{ew} -algebra is k -potent for some k .

COROLLARY 4.2 *If a subdirectly irreducible FL_{ew} -algebra M satisfies E_k for some k , then it has the unique coatom.*

Proof: Suppose that M satisfies E_k . Since M is subdirectly irreducible, there exists an element $a (< 1)$ such that for each $x < 1$ there exists a positive integer m such that $x^m \leq a$ holds. Then $x^k \leq a$ must hold also for each $x < 1$. Thus, by Lemma 4.2 M has the unique coatom. \square

A related result for BCK-algebras was shown in [57]. From the above lemma, it follows immediately that every finite subdirectly irreducible FL_{ew} -algebra has a single coatom. It is easy to see that a subdirectly irreducible FL_{ew} -algebras with the unique coatom does not always satisfy E_k for some k . In the following, we will give an interesting example of a subdirectly irreducible FL_{ew} -algebra with no coatoms. This special FL_{ew} -algebra has been discussed by Pavelka [58], Hájek, Godo and Esteva [31] etc. in connection with fuzzy logic.

Let M be the interval $[0, 1]$, i.e. the set of all real numbers between 0 and 1. Then M forms a bounded lattice with the natural order \leq . Moreover, $\langle M, \times, 1 \rangle$ is a commutative monoid satisfying $x \times y \leq x$, where \times denotes the usual multiplication. Since M is linearly ordered, the distributivity of \times with \vee follows from the above inequality. Now, define a binary operation \rightarrow on M by

$$y \rightarrow z = \begin{cases} z/y & \text{if } y > z, \\ 1 & \text{otherwise.} \end{cases}$$

Then, we can show that $x \times y \leq z$ if and only if $x \leq y \rightarrow z$. Thus, M forms an FL_{ew} -algebra. Clearly, M has no coatoms. Now, let α be an arbitrary real number such that $0 < \alpha < 1$. Then, we can show that for each $r < 1$ there exists a positive number m such that $r^m \leq \alpha$. This means that M is subdirectly irreducible.

The logic determined by the above FL_{ew} -algebra is called *product logic*, whose axiomatization is given in [31]. A class of FL_{ew} -algebras related to product logic, called *product algebras*, is discussed in [1, 30, 31] etc..⁹

The following fact, remarked in [3], should be pointed out here. Suppose that α is a real number such that $0 < \alpha < 1$. Define an operation \circ on the interval $[\alpha, 1]$ by

$$x \circ y = \max\{\alpha, x \times y\}.$$

Then, this with the operation \rightarrow satisfying

$$y \rightarrow z = z/y \text{ if } y > z, \text{ and } = 1 \text{ otherwise.}$$

determines an FL_{ew} -algebra, which we call I_α . Obviously, this is also a subdirectly irreducible FL_{ew} -algebra without coatoms. The logic determined by this FL_{ew} -algebra is equal to Łukasiewicz's infinitely many-valued logic. In fact, the mapping ϕ from $[0, 1]$ to $[\alpha, 1]$ defined by $\phi(x) = \alpha^{1-x}$ is a monotone increasing map satisfying $\phi(0) = \alpha$, $\phi(1) = 1$ and $\phi(\max\{0, x + y - 1\}) = \phi(x) \circ \phi(y)$, and hence it is an isomorphism (see also [30]).

In the dual form, this fact can be restated as follows. Let R be the set of all real numbers as an additive ordered Abelian group. Then the mapping ψ from the MV-algebra $R[1]$ (see Chang [15] for the definition) to I_α by $\psi(x) = \alpha^x$ is an isomorphism.

Let Lin be the formula $(p \rightarrow q) \vee (q \rightarrow p)$, which is sometimes called *the axiom of prelinearity*. It is called also *the (algebraic) strong de Morgan law* in [38] and [59]. Using Theorem 4.1 we can show the following (see also [33] Theorem 4.8).

LEMMA 4.3 *For any subdirectly irreducible FL_{ew} -algebra M , the formula Lin is valid in M if and only if M is linearly ordered.*

Proof: The if-part is trivial. Suppose that Lin is valid in M . This implies that $(a \rightarrow b) \vee (b \rightarrow a) = 1$ for all $a, b \in M$. By Theorem 4.1, either $a \rightarrow b = 1$ or $b \rightarrow a = 1$ holds. Thus, either $a \leq b$ or $b \leq a$. Hence, M is linearly ordered.

COROLLARY 4.3 *If the formula Lin is valid in an FL_{ew} -algebra M then M is a distributive lattice.*

Proof: Suppose that M is subdirectly represented by the set of subdirectly irreducible FL_{ew} -algebras N_j for $j \in J$. Then, Lin is valid in any N_j . Therefore,

⁹(†) There are considerable developments in the study of product logic after 2001.

by Lemma 4.3 each N_j is linearly ordered and hence is distributive. In other words, the formula $\text{Dis} : (p \wedge (q \vee r)) \rightarrow (p \wedge q) \vee (p \wedge r)$ is valid in each of them. Thus, Dis is valid in M and hence M is distributive. \square

As for a direct proof of the above corollary, see Corollary 1 of Proposition 4.3.4 in [59], or Lemma 2.4 (3) in [33].

Suppose that M and N are finite subdirectly irreducible FL_{ew} -algebras. It is interesting to see when they determine the same logics. For this purpose, we will show a theorem in the following which is obtained by modifying the proof of Jankov's theorem for Heyting algebras (see [37]).

Let M be an arbitrary finite, subdirectly irreducible FL_{ew} -algebra and ω be the coatom of M . For each element $x \in M$, take a propositional variable p_x in such a way that p_y and p_z are distinct whenever $y \neq z$. Define a set of formulas Δ_M by

$$\Delta_M = \{(p_x \wedge p_y) \equiv p_{x \wedge y}, (p_x \vee p_y) \equiv p_{x \vee y}, (p_x \cdot p_y) \equiv p_{x \cdot y}, \\ (p_x \rightarrow p_y) \equiv p_{x \rightarrow y}, (\neg p_x) \equiv p_{\sim x} : x, y \in M\}$$

Let $(\Delta_M)^\dagger$ be the conjunction of all formulas in Δ_M . For each positive integer k , define the formula $X_M(k)$, called the *Jankov formula of order k for M* , by $X_M(k) = ((\Delta_M)^\dagger)^k \rightarrow p_\omega$.¹⁰ Then, we have the following Jankov's theorem for FL_{ew} -algebras.

THEOREM 4.2 *Let M be a finite, subdirectly irreducible FL_{ew} -algebra, and let N be any FL_{ew} -algebra satisfying E_k . Then, $X_M(k)$ is not valid in N if and only if M is embeddable into a quotient algebra of N .*

Proof: Suppose that h is an embedding of M into a quotient algebra N/F of N . Define a valuation v on N/F by $v(p_x) = h(x)$ for each $x \in M$. Then, it is clear that $v(X_M(k)) = 1 \rightarrow h(\omega) = h(\omega) < 1$. Thus, $X_M(k)$ is not valid in N/F , and nor is it in N . Conversely, suppose that the formula $X_M(k)$ becomes false under a valuation w on N . Let $w((\Delta_M)^\dagger) = d$ and $w(p_\omega) = e$. By our assumption, $d^k \not\leq e$. Since N satisfies E_k , the set G defined by $G = \{x : d^k \leq x\}$ forms a filter such that $e \notin G$. Now, consider the quotient algebra N/G of N and its valuation w^* defined by $w^*(q) = \|w(q)\|$ for every propositional variable q . Here $\|u\|$ denotes the equivalence class induced by G , to which u belongs. Then $w^*((\Delta_M)^\dagger) = 1$ and $w^*(p_\omega) = \|e\| < 1$. From this, it follows that the mapping $g : M \rightarrow N/G$, defined by $g(x) = w^*(p_x)$ for each $x \in M$, is an embedding. \square

By assuming moreover that M satisfies also E_k , we can give another condition which is equivalent to either of the conditions in the above theorem.

¹⁰In [45], these Jankov formulas are used in the study of *splittings* in the variety of FL_{ew} -algebras.

COROLLARY 4.4 *Suppose that M is a finite, subdirectly irreducible FL_{ew} -algebra and N is an arbitrary FL_{ew} -algebra. When both satisfy E_k , M is embeddable into a quotient algebra of N if and only if $L(N) \subseteq L(M)$.*

Proof: The only-if part is trivial. Suppose that $L(N) \subseteq L(M)$. Since M satisfies E_k , $\chi_M(k)$ is not valid in M . Hence it is not valid in N , either. Then, by Theorem 4.2, there is an embedding of M into a quotient algebra of N . \square

A Heyting algebra is an FL_{ew} -algebra satisfying E_1 , i.e. $x^2 = x$. Theorem 4.2 with Corollary 4.4 gives us the original form of Jankov's theorem for Heyting algebras.

Let E_k be the formula $p^k \rightarrow p^{k+1}$ for any $k > 0$, where A^j denotes the formula $A \cdot \dots \cdot A$ with j times A . Obviously, the formula E_k is a syntactic expression of E_k .

COROLLARY 4.5 *Suppose that both M and N are finite subdirectly irreducible FL_{ew} -algebras. If the logic $L(M)$ is equal to $L(N)$ then M is isomorphic to N .*

Proof: Suppose that $L(M)$ is equal to $L(N)$. Since M is finite, we can assume that E_k is in $L(M)$ (and hence in $L(N)$) for some k . By Corollary 4.4, there is an embedding h of M into a quotient algebra of N . Hence, $|M| \leq |N|$, where $|S|$ denotes the cardinality of a set S , and M and N are universes of algebras M and N , respectively. Similarly, since there is an embedding of N into a quotient algebra of M , $|N| \leq |M|$, and hence $|M| = |N|$. Thus h must be an isomorphism of M onto N .

5 SOME AXIOMATIZABLE EXTENSIONS OF FL_{ew}

To give an overview of the structure of all logics over FL_{ew} , we will introduce some basic logics among them, which are finitely axiomatizable. Let us consider the following axioms:

EM: $p \vee \neg p$ (excluded middle)

DN: $\neg\neg p \rightarrow p$ (double negation)

Con: $(p \rightarrow (p \rightarrow q)) \rightarrow (p \rightarrow q)$ (contraction)

WCon: $(p \rightarrow \neg p) \rightarrow \neg p$ (weak contraction)

P: $((p \rightarrow q) \rightarrow p) \rightarrow p$ (Peirce's law)

WP: $(\neg p \rightarrow p) \rightarrow p$ (weak Peirce's law)

Lin: $(p \rightarrow q) \vee (q \rightarrow p)$ (prelinearity)

Dis: $(p \wedge (q \vee r)) \rightarrow ((p \wedge q) \vee (p \wedge r))$ (distributive law)

Note that $WCon$ and WP are obtained from Con and P , respectively, by replacing q by \perp . The axiom Con will play essentially the same role as the contraction rule (and also as the formula E_1 introduced in the previous section). Therefore $FL_{ew}[Con]$ is equal to the intuitionistic logic Int , which is equal also to FL_{ew} with the axiom $p \rightarrow p^2$. On the other hand, $WCon$ is equivalent to $\neg p^2 \rightarrow \neg p$, which is the contrapositive of $p \rightarrow p^2$, and $FL_{ew}[WCon]$ is equal to FL_{ew} with the axiom $\neg(p \wedge \neg p)$. It is clear that $FL_{ew}[DN]$ is equal to the logic without the contraction rule studied by Grišin.

LEMMA 5.1 *Each of the logics $FL_{ew}[WP]$, $FL_{ew}[P]$, $FL_{ew}[EM]$ and $FL_{ew}[WCon, DN]$ is equal to classical logic Cl .*

Proof: As mentioned above, any instance of WP is provable in $FL_{ew}[P]$. It is obvious that any instance of P is provable in Cl . To show that each formula A provable in Cl is also provable in $FL_{ew}[EM]$, it is enough to show that Con is provable in $FL_{ew}[EM]$, since $FL_{ew}[Con]$ is equal to Int and $Int[EM]$ is equal to Cl . We note that the formula $(\neg r \vee s) \equiv (r \rightarrow s)$ is provable in $FL_{ew}[EM]$, using the fact that

$$(r \vee \neg r) \rightarrow ((r \rightarrow s) \rightarrow (\neg r \vee s))$$

is provable in FL_{ew} . Thus,

$$((p \rightarrow (p \rightarrow q)) \rightarrow (p \rightarrow q)) \equiv ((\neg p \vee \neg p \vee q) \rightarrow (\neg p \vee q))$$

is provable. But, the righthand side of the above equivalence is always provable. Hence, Con is provable in $FL_{ew}[EM]$. To show that EM is provable in $FL_{ew}[WP]$, let Q be $p \vee \neg p$. Now, since $\neg Q \rightarrow Q$ is provable in FL_{ew} and $(\neg Q \rightarrow Q) \rightarrow Q$ is an instance of WP , we have that Q , i.e. $p \vee \neg p$, is provable in $FL_{ew}[WP]$.

Lastly, we will show that $FL_{ew}[WCon, DN]$ is equal to Cl . Since $FL_{ew}[Con]$ is equal to Int , it is enough to show that $p \rightarrow p^2$ is provable in $FL_{ew}[WCon, DN]$. By $WCon$, $\neg p^2 \rightarrow \neg p$ is provable. By taking the contraposition and using DN , $p \rightarrow p^2$ follows. This completes our proof. \square

Note that $FL_e[P]$ is in fact equal to Cl as the weakening rule is derivable in it. (This is due to K. Fujita.) A part of the above lemma is essentially obtained in Theorem 7.31 of [63], which says that complemented FL_{ew} -algebras are exactly same as Boolean algebras. It can be shown (by using e.g. the cut elimination theorem) that Grišin's logic $FL_{ew}[DN]$ is strictly weaker than the classical logic, and hence EM is not provable in $FL_{ew}[DN]$. Also, we can show that EM is provable in $FL_e[WCon, DN]$. (On the other hand, probably DN will not be provable in $FL_{ec}[EM]$.)

Figure 5.1 shows the inclusion relationship between logics mentioned in the above. In the following, we will show that all of the inclusions described in the figure are proper.

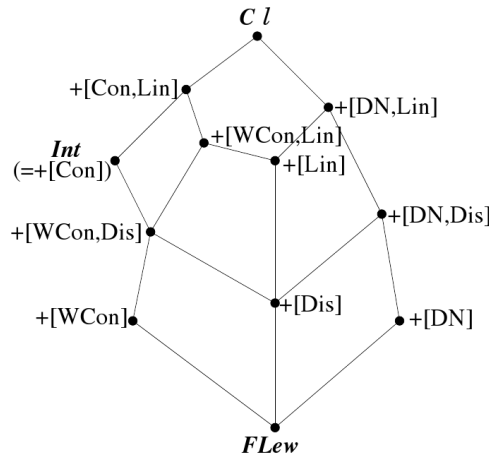


Figure 5.1.

(1) The logic $FL_{ew}[Con, Lin]$, i.e. $Int[Lin]$, is Gödel logic, which is strictly weaker than Cl . Thus, $FL_{ew}[Lin]$, $FL_{ew}[Dis]$, and FL_{ew} are strictly weaker than $FL_{ew}[DN, Lin]$, $FL_{ew}[DN, Dis]$ and $FL_{ew}[DN]$, respectively.

(2) Gödel logic $Int[Lin]$ is strictly stronger than Int . Thus, logics $FL_{ew}[WCon, Lin]$ and $FL_{ew}[Lin]$ are strictly stronger than $FL_{ew}[WCon, Dis]$ and $FL_{ew}[Dis]$, respectively.

We will show that Lin is not provable in $FL_{ew}[DN, Dis]$. Let U_1 be the FL_{ew} -algebra with the universe $\{0, d, c, b, a, 1\}$ such that $0 < d < b < a < 1$, $0 < d < c < a < 1$ but b is incomparable with c . Moreover, assume that $a^2 = a, b^2 = b, c^2 = c, a \cdot b = b, a \cdot c = c, a \cdot d = 0$ (thus, $x \cdot d = 0$ if $x < a$) and $b \cdot c = 0$. Then, it can be seen that U_1 is in fact a subdirectly irreducible, distributive FL_{ew} -algebra satisfying $\sim\sim x = x$ for any x . Clearly, it is not linearly ordered. Thus, Lin is not valid in it.

(3) To see that $FL_{ew}[DN, Lin]$ is strictly weaker than Cl , consider the FL_{ew} -algebra U_2 with the universe $\{0, a, 1\}$ such that $0 < a < 1$ and $a^2 = 0$. Since $\sim a = a$ holds, DN is valid but EM is not.

From this it follows that $FL_{ew}[Lin]$ is strictly weaker than $FL_{ew}[WCon, Lin]$. For, if $WCon$ were provable in $FL_{ew}[Lin]$ then it would be provable in $FL_{ew}[DN, Lin]$. Then, EM should be derived in $FL_{ew}[DN, Lin]$, since by Lemma 5.1 it is provable in $FL_{ew}[WCon, DN]$. This is a contradiction. Similarly, we can show that $FL_{ew}[Dis]$ and FL_{ew} are strictly weaker than $FL_{ew}[WCon, Dis]$ and $FL_{ew}[WCon]$, respectively.

(4) We will show next that Con is not provable in $FL_{ew}[WCon, Lin]$. Let U_3 be the FL_{ew} -algebra with the universe $\{0, b, a, 1\}$ such that $0 < b < a < 1$ and

$b = a^2 = a^3$. It is easy to see that $a \rightarrow b = a$ and $\sim a = \sim b = 0$. It is clear that $WCon$ is valid in a given FL_{ew} -algebra if and only if $\sim x^2 = \sim x$ for any x in it. So, $WCon$ is valid in U_3 . On the other hand, Con is not valid in it, as $a \rightarrow a^2 = a < 1$. As a consequence, Con is not provable in $FL_{ew}[WCon]$, either.

(5) In [29], Grišin introduced a cut-free system for $FL_{ew}[DN]$. Using this system, we can easily show that Dis is not provable in $FL_{ew}[DN]$, and *a fortiori* is not provable in FL_{ew} . Now we show that Dis is not provable in $FL_{ew}[WCon]$. Let U_4 be the FL_{ew} -algebra with the universe $\{0, d, c, b, a, 1\}$ such that $0 < b < a < 1, 0 < d < c < a < 1$ and both c and d are incomparable with b . (Thus, U_4 can be obtained from the nonmodular (and hence nondistributive) lattice N_5 , called the *pentagon*, by adding the new greatest element 1 . See e.g. [19] Chapter 6.) Define also that $a^2 = a, a \cdot b = b^2 = b, a \cdot c = a \cdot d = c^2 = c \cdot d = d^2 = d$ and $b \cdot c = b \cdot d = 0$. Then, it can be seen that U_4 is a subdirectly irreducible, nondistributive FL_{ew} -algebra. On the other hand, we can show that $\sim a = 0, \sim b = c$ and $\sim c = \sim d = b$. Using this, we can show that $WCon$ is valid in U_4 . Thus, Dis is not provable in $FL_{ew}[WCon]$.

6 SIMPLE AND SEMISIMPLE FL_{ew} -ALGEBRAS

In this section, we will discuss simple and semisimple FL_{ew} -algebras. As usual, a *simple* FL_{ew} -algebra M is defined to be a non-degenerate FL_{ew} -algebra which has only two filters $\{1\}$ and M itself. It is easy to see that for any filter F of a given FL_{ew} -algebra M^* the quotient algebra M^*/F is simple if and only if F is a maximal filter. Similarly to the case of Wajsberg algebras, the following holds. (See [14] Theorem 4.7 and [24] Corollary 1 of Theorem 17.)

LEMMA 6.1 *An FL_{ew} -algebra M is simple if and only if for any $x (< 1)$ in M there exists a positive integer m such that $x^m = 0$.*

Clearly, any simple FL_{ew} -algebra is subdirectly irreducible. It is well-known that for a Heyting algebra M , if it is simple then it is just the two-valued Boolean algebra B_2 and therefore satisfies that $\sim \sim x = x$ for any x . On the other hand, this equation does not hold always in simple FL_{ew} -algebras. Also, we know that any simple Wajsberg algebra is linearly ordered (see e.g. [14, 41, 24]). But, this does not hold either for FL_{ew} -algebras, in general. In fact, we can give an example of a simple FL_{ew} -algebra N which is not distributive (as a lattice) and hence is not linearly ordered. Let M_3 be a nondistributive lattice, called the *diamond* (see e.g. [19]). Let ω and 0 be the greatest element and the least element of M_3 . Then N is obtained from M_3 by adding a new element 1 and requiring that $\omega < 1$. Define operations \cdot and \rightarrow on N as follows; $x \cdot y = 0$ for $x, y \in M_3$ and $x \cdot 1 = 1 \cdot x = x$ for any $x \in N$, and $x \rightarrow y = 1$ if $x \leq y, = y$ if $x = 1 > y$ and $= \omega$ otherwise. Then, it is easily checked that N is a nondistributive, simple FL_{ew} -algebra. \square

Any FL_{ew} -algebra I_a for $0 < a < 1$ introduced in Section 4 is simple but does not have any coatom. Each Łukasiewicz's $n + 1$ -valued model L_{n+1} is a simple FL_{ew} -algebra, which is isomorphically represented as the FL_{ew} -algebra with the set $\{1, a, a^2, \dots, a^n\}$ such that $a^i > a^{i+1}$ for each $i = 1, \dots, n - 1$ and $a^n = 0$. For the simplicity's sake, let us denote 1 by a^0 . Then the operation \rightarrow on this FL_{ew} -algebra can be defined by $a^k \rightarrow a^m = a^{\max\{0, m-k\}}$.

This representation might be convenient when we compare Łukasiewicz's models with other MV-algebras. For example, let us take the MV-algebra C discussed by Chang in [14], which is call S_1^ω in Komori [41]. Then, C can be isomorphically represented as the residuated lattice with the set $\{1, a, a^2, \dots, a^n, \dots, \dots, \sim a^n, \dots, \sim a^2, \sim a, 0\}$ such that $a^{i-1} > a^i, a^i > \sim a^j$ and $\sim a^j > \sim a^{j-1}$ for each positive integer i, j . The operations \cdot and \rightarrow are defined as follows:

- $a^k \cdot a^m = a^{k+m},$
- $a^k \cdot (\sim a^m) = \sim a^{\max\{m-k, 0\}},$
- $(\sim a^k) \cdot (\sim a^m) = 0,$
- $a^k \rightarrow a^m = a^{\max\{0, m-k\}},$
- $(\sim a^k) \rightarrow a^m = 1,$
- $a^k \rightarrow (\sim a^m) = \sim a^{k+m},$
- $(\sim a^k) \rightarrow (\sim a^m) = a^{\max\{k-m, 0\}}.$

It is clear that C is not simple and has only two proper filters $\{1\}$ and $\{a^n : n \in \mathbb{N}\}$. □

For an FL_{ew} -algebra M , let Φ_M be the set of all maximal filters of M . We note here that a filter F is maximal if and only if for any $u \in M$ either $u \in F$ or $\sim u^k \in F$ for some $k \geq 1$. Define the *radical* Rad_M of M by $Rad_M = \bigcap_{F \in \Phi_M} F$. An FL_{ew} -algebra M is *semisimple* if M can be represented by a subdirect product of simple FL_{ew} -algebras, or equivalently if $Rad_M = \{1\}$ (see e.g. [48]). The proof of the following theorem is essentially due to Grišin [28]. In [29] he used the result to show that every *free* involutive FL_{ew} -algebra is semisimple. From this, it follows that the logic $FL_{ew}[DN]$ can be characterized by the class of simple involutive FL_{ew} -algebras. Now, for any x, y in a given FL_{ew} -algebra, define $x + y$ by $x + y = \sim(\sim x \cdot \sim y)$. We can show easily that the operation $+$ is associative and moreover that $(x + y) + z = \sim(\sim x \cdot \sim y \cdot \sim z)$. For an element x and a positive integer m , $\tilde{m}x$ denotes $\sim(\sim x)^m$. Then $\tilde{1}x = \sim\sim x$ and $\tilde{m}x = x + \dots + x$ with m times x when $m > 1$.

THEOREM 6.1 *For any x in a given FL_{ew} -algebra M , $x \in Rad_M$ if and only if for any $n \geq 1$ there exists $m \geq 1$ such that $\tilde{m}(x^n) = 1$.*

Proof: Suppose first that for any $n \geq 1$ there exists $m \geq 1$ such that $\tilde{m}(x^n) = 1$. Suppose that $x \notin \text{Rad}_M$. Then there exists a maximal filter F such that $x \notin F$. Since F is maximal, there exists a $k \geq 1$ such that $\sim x^k \in F$. There exists $m \geq 1$ such that $\tilde{m}(x^k) = 1$, i.e., $(\sim x^k)^m = 0$. Thus, $0 \in F$, which contradicts the fact that F is proper.

Conversely, suppose that there exists $n (\geq 1)$ such that $\tilde{m}(x^n) \neq 1$ for any m . If $(\sim(x^n))^m = 0$ then $\tilde{m}(x^n) = \sim(\sim(x^n))^m = 1$, which is a contradiction. Thus, $(\sim(x^n))^m > 0$ for any m . Let $z = \sim(x^n)$ and H be the filter generated by z . Clearly, H is proper as $z^m > 0$ for any m . By Zorn's lemma, there exists a maximal filter G such that $H \subseteq G$. Now, suppose that $x \in G$. Then x^n must be also in G . But this is a contradiction, since $z = \sim(x^n) \in G$. Hence, $x \notin G$ and therefore it does not belong to Rad_M . \square

This result was also shown in Höhle [33]. As a corollary of Theorem 6.1, we have the following result on Wajsberg algebras, which was shown by Font, Rodríguez and Torrens in [24]. (See also [33].)

COROLLARY 6.1 *For any x in a given Wajsberg algebra M , $x \in \text{Rad}_M$ if and only if for any $n \geq 1$, $\sim(x^n) \leq x$.*

Proof: By Theorem 6.1, it suffices to show that in any Wajsberg algebra M , $\sim(x^n) \leq x$ for any $n \geq 1$ if and only if for any $n \geq 1$ there exists $m \geq 1$ such that $\tilde{m}(x^n) = 1$.

Suppose first that $\sim(x^n) \leq x$. We will show that $\widetilde{n+1}(x^n) = 1$. From the assumption, we have $(\sim(x^n))^n \leq x^n$. Thus, $\sim x^n \leq \sim(\sim(x^n))^n = \tilde{n}(x^n)$ and hence $1 = \sim x^n \rightarrow n(x^n) = x^n + \tilde{n}(x^n) = \widetilde{n+1}(x^n)$. Conversely, suppose that $\sim(x^n) \not\leq x$ for some $n \geq 1$. Take any subdirect representation $f : M \rightarrow \prod_i M_i$ with subdirectly irreducible factors M_i . Note that each M_i is linearly ordered by Lemma 4.3. Let p_j be the j -th projection function of the direct product $\prod_i M_i$, and let $z_j = (p_j \circ f)(x)$. By our assumption, $\sim(z_j)^n \not\leq z_j$ for some j . Since M_j is linearly ordered, $z_j \leq \sim(z_j)^n$ holds. Therefore $(z_j)^{n+1} = 0$. Then, $\tilde{m}(x^{n+1}) = 1$ never hold for any m . This completes the proof of our lemma. \square

The next theorem gives us a sufficient condition for a FL_{e_w} -algebra to be semisimple. Here, we will introduce the following condition on FL_{e_w} -algebras:

$$EM_k : x \vee (\sim x^k) = 1 \text{ for any } x.$$

Note that EM_1 corresponds to an algebraic form of the excluded middle EM.

THEOREM 6.2 *The following three conditions are mutually equivalent for any subdirectly irreducible FL_{e_w} -algebra M .*

1. M is simple and satisfies E_k ,
2. for each $x \in M$, if $x < 1$ then $x^k = 0$,

3. M satisfies EM_k .

Proof: Assume that M is simple and satisfies EM_k . If $x < 1$ then $1 > x \geq x^2 \geq \dots \geq x^k = x^{k+1}$. The filter F generated by x is expressed as $\{z \in M : x^k \leq z\}$. If $x^k > 0$ then F is neither equal to $\{1\}$ nor equal to M . But this contradicts our assumption that M is simple.

Next, suppose that $x < 1$ implies $x^k = 0$ for any $x \in M$. Take any $y \in M$. If $y = 1$ then clearly $y \vee (\sim y^k) = 1$. If $y < 1$, $y^k = 0$ by our assumption. Then, $\sim y^k = 1$ and hence $y \vee (\sim y^k) = 1$ holds also in this case.

Finally, suppose that M satisfies EM_k . Suppose that $x < 1$. Since $x \vee (\sim x^k) = 1$, $\sim x^k = 1$ by Theorem 4.1. Then, $x^k = x^k \cdot 1 = x^k \cdot \sim x^k = 0$. Thus, M is simple. Next, it is easy to see that for each y , $y^k \cdot y \leq y^{k+1}$ and $y^k \cdot \sim y^k = 0 \leq y^{k+1}$. Therefore, $y^k \cdot (y \vee \sim y^k) = (y^k \cdot y) \vee (y^k \cdot \sim y^k) \leq y^{k+1}$. But, $y \vee \sim y^k = 1$ by our assumption. Thus, $y^k \leq y^{k+1}$. Since the converse inequality holds always, we have that $y^k = y^{k+1}$ for any y . \square

Let us take τ for k in the above theorem. Then, we have the following well-known result: For any subdirectly irreducible FL_{ew} -algebra M , the following three conditions are mutually equivalent (cf. Lemma 5.1).

1. M is simple and satisfies $x^2 = x$ for any x , i.e., M is a simple Heyting algebra,
2. M is isomorphic to the two-valued Boolean algebra B_2 ,
3. M satisfies $x \vee \sim x = 1$ for any x .

A characterization of the variety defined by the condition EM_k and of the variety defined by EM_k from universal algebra is given in [43]. Following Chang [14] and Komori [41], for any FL_{ew} -algebra M , define the *order* $o(M)$ of M as follows: if k is the smallest number j such that M satisfies EM_j , let $o(M) = k$, and if there exists no such k , let $o(M) = \omega$. For each k , let EM_k be the formula $p \vee \sim p^k$, which corresponds to the identity EM_k . It is obvious that for any FL_{ew} -algebra M , $o(M) = k$ for a finite k if and only if k is the smallest number j such that $FL_{ew}[EM_j] \subseteq L(M)$.

COROLLARY 6.2 *Any FL_{ew} -algebra with a finite order is semisimple.*

THEOREM 6.3 *The intersection of $FL_{ew}[E_n]$ for $n < \omega$ is equal to FL_{ew} .*

Proof: Our theorem follows from the finite model property of FL_{ew} proved by Okada and Terui in [50]. Suppose that a formula A is not provable in FL_{ew} . Then by the finite model property of FL_{ew} , A is not valid in a finite FL_{ew} -algebra M . Suppose that M satisfies E_k . Then, $L(M)$ is a logic over $FL_{ew}[E_k]$ and thus, A is not provable in $FL_{ew}[E_k]$. \square

In Kowalski and Ono [44], it is proved that the logic FL_{ew} is determined by the class of *finite, simple* residuated lattices. This improves the result by Okada and Terui on the finite model property, and moreover implies immediately the following, using Theorem 6.2.

THEOREM 6.4 *The intersection of $FL_{ew}[EM_n]$ for $n < \omega$ is equal to FL_{ew} .*

Using the sequence $\{E_k\}_k$, we introduce a classification $\{W_k\}_k$ of logics over FL_{ew} in the following way.

- $W_1 = \{L : FL_{ew}[E_1] \subseteq L \subseteq Cl\}$,
- $W_{n+1} = \{L : FL_{ew}[E_{n+1}] \subseteq L \subseteq Cl\} - W_n$,
- $W_\omega = \{L : FL_{ew} \subseteq L \subseteq Cl\} - \bigcup_{n < \omega} W_n$.

Clearly, the class W_1 is exactly the class of all intermediate logics. This classification is *orthogonal* to the class of Łukasiewicz's many-valued logics in the sense that for each $n > 0$ only the logic $L(\mathfrak{L}_{n+1})$ belongs to W_n and that $L(\mathfrak{L})$ belongs to W_ω . By the definition of the classification, each W_n has the smallest element $FL_{ew}[E_n]$ for each $0 < n \leq \omega$, where $FL_{ew}[E_\omega]$ means FL_{ew} . On the other hand, W_n has no greatest element in general, as shown in the next section.

7 NEIGHBORS OF CLASSICAL LOGIC

A logic L over FL_{ew} is an *immediate predecessor* of classical logic Cl if (1) $L \subset Cl$ and (2) $L \subseteq L' \subset Cl$ implies $L' = L$ for any logic L' . In this section, we will study immediate predecessors of Cl . This, we hope, will be the first step in studying the whole structure of the lattice consisting of logics over FL_{ew} .¹¹ The set of all immediate predecessors of Cl is denoted by $IP(Cl)$ in the following.

Suppose that $L \in IP(Cl)$ and that L is characterized by an FL_{ew} -algebra M . Let us assume moreover that M is subdirectly represented by $\{N_j\}_{j \in J}$. Then, $L(M) \subseteq L(N_j) \subseteq Cl$ for each $j \in J$. Since $L(M)$ is an immediate predecessor of Cl , some $L(N_j)$ must be equal to $L(M)$. Thus, we can assume that L is characterized by a *subdirectly irreducible* FL_{ew} -algebra M from the beginning. \square

We have already seen some of immediate predecessors of Cl . If we restrict our attention to logics over the intuitionistic logic Int , i.e. intermediate logics, then there exists the single immediate predecessor of Cl , which is determined by the three valued Heyting algebra H_3 . In other words, the class W_1 has the single immediate predecessor H_3 of Cl . On the other hand, Komori [41] discussed the lattice structure of logics over $L(\mathfrak{L})$ and proved that in this class L is an immediate predecessor of Cl if and only if either $L = L(\mathfrak{L}_{n+1})$ with a prime number n or $L = L(C)$, where C is the MV-algebra of Chang (see the previous section). For each n such that $0 < n \leq \omega$, let W_n^* be the class of maximal elements in W_n , and W_n^{IP} be the set $W_n \cap IP(Cl)$, i.e. the set of all immediate predecessors of Cl in W_n . Clearly, W_n^{IP} is a subset of W_n^* when $n > 1$.

¹¹(†) For recent developments of the study in this direction, see Chapter 9 of [25].

THEOREM 7.1 *For each $n < \omega$, $L(\mathfrak{L}_{n+1})$ belongs to W_n^* . Also, $L(C)$ belongs to W_ω^* .*

Proof: Suppose that $L(\mathfrak{L}_{n+1}) \subseteq L$ and $L \in W_n$. Clearly, L is a logic over $L(\mathfrak{L})$. Using results in Komori [41], there exists a finite set I of natural numbers such that $L = \bigcap_{i \in I} L(\mathfrak{L}_{i+1})$. Then, by Lindenbaum's result, each $i \in I$ must be a divisor of n . If $i < n$ for all $i \in I$, then L must be in W_j where $j = \max\{i : i \in I\} < n$. This is a contradiction. Thus, $i = n$ for some $i \in I$. Then, $L \subseteq L(\mathfrak{L}_{n+1})$ and hence $L = L(\mathfrak{L}_{n+1})$. Therefore, $L(\mathfrak{L}_{n+1})$ is maximal in W_n . Since $L(C)$ belongs to W_ω^{IP} , the second part is obvious. \square

We will show in the rest of this section that there are many other immediate predecessors of Cl . Note that each \mathfrak{L}_{n+1} is simple, but neither H_3 nor C is so. It is clear that every FL_{ew} -algebra M contains M itself and $\{0, 1\}$ as its subalgebras. The latter is obviously isomorphic to the two-valued Boolean algebra B_2 . A subalgebra K of an FL_{ew} -algebra M is *trivial* if K is either M or $\{0, 1\}$. The following theorem gives us a condition for a *finite* subdirectly irreducible FL_{ew} -algebra to determine an immediate predecessor of Cl .

THEOREM 7.2 *Suppose that M is a finite, subdirectly irreducible (non-degenerate) FL_{ew} -algebra. Then, $L(M)$ is either equal to classical logic Cl or an immediate predecessor of Cl if and only if*

1. *every subalgebra of M is trivial,*
2. *for any proper filter F of M , either $F = \{1\}$ or the quotient algebra M/F is isomorphic to B_2 .*

To show Theorem 7.2, we use Jónsson's lemma [39] on congruence-distributive variety (see also [5]). Recall that the class of FL_{ew} -algebras forms a congruence-distributive variety by Propositions 3.4 and 3.5. In the following, $V(\mathcal{K})$ and $V(\mathcal{K})_{SI}$ mean the variety generated by a class \mathcal{K} of algebras and the class of subdirectly irreducible algebras in $V(\mathcal{K})$, respectively. Also, $H(\mathcal{K})$, $S(\mathcal{K})$ and $P_U(\mathcal{K})$ denote the class of homomorphic images of members of \mathcal{K} , the class of subalgebras of members of \mathcal{K} , and the class of ultraproducts of nonempty families of members of \mathcal{K} , respectively.¹² We can state Jónsson's lemma applied to the present case as follows.

PROPOSITION 7.1 *Let \mathcal{K} be a class of FL_{ew} -algebras. Then, $V(\mathcal{K})_{SI} \subseteq HSP_U(\mathcal{K})$. In particular, if M_1, \dots, M_n are finite FL_{ew} -algebras, then $V(\{M_1, \dots, M_n\})_{SI} \subseteq HS(\{M_1, \dots, M_n\})$.*

PROOF OF THEOREM 7.2. Only-if part is trivial. We assume that a finite subdirectly irreducible M satisfies both conditions 1 and 2. Moreover, suppose that $L(M) \subseteq L \subset Cl$ for a logic L . The logic L can be represented as $\bigcap_{i \in I} L(N_i)$ for some non-degenerate, subdirectly irreducible FL_{ew} -algebras N_i for $i \in I$. Since $L(M) \subseteq L(N_i)$, $N_i \in V(\{M\})_{SI} \subseteq HS(\{M\})$, by using Proposition 7.1. By

¹²(†) For more information on universal algebra, see [11].

condition 1, $S(\{M\}) = \{M, B_2\}$. Thus, each N_i must be the homomorphic image of some $K \in \{M, B_2\}$ by a homomorphism h . Hence, the kernel $Ker(h)$ of h is a proper filter of K , and $N_i \cong K/Ker(h)$, by the homomorphism theorem. By condition 2, either $Ker(h) = \{1\}$ or the quotient algebra N_i is isomorphic to B_2 . In other words, N_i is isomorphic either to $K \in \{M, B_2\}$ or to B_2 . Since $L = \bigcap_{i \in I} L(N_i) \subset Cl$, N_j must be isomorphic to M for some $j \in I$. Therefore, $L = L(M)$ as $L(M) \subseteq L(B_2)$. Thus, $L(M)$ is an immediate predecessor of Cl . \square

COROLLARY 7.1 *If M is a finite, simple (non-degenerate) FL_{ew} -algebra without non-trivial subalgebras, then $L(M)$ is equal either to classical logic Cl or to an immediate predecessor of Cl .*

Proof: If M is simple, it has only $\{1\}$ as its proper filter. Thus, condition 2 in Theorem 7.2 is satisfied. \square

This corollary gives us a proof of a well-known result which says that $L(\mathcal{L}_{n+1})$ is an immediate predecessor of Cl when n is prime. Note that in [61] K. Swirydowicz showed the existence of infinitely many immediate predecessors of Cl over the classical linear logic, i.e. $FL_e[DN]$, by using the similar criterion as Corollary 7.1.

By Theorem 6.2 and Corollary 6.2, we have the following.

COROLLARY 7.2 *Suppose that an immediate predecessor L of Cl is determined by an FL_{ew} -algebra M . Then, M is simple and k -potent if and only if L is a logic over $FL_{ew}[EM_k]$.*

By using Corollary 7.1, we will give here an example of simple, linearly ordered FL_{ew} -algebra K which determines an immediate predecessor of Cl , but which is not isomorphic to any of Łukasiewicz's models. The universe of K consists of 6 elements, which are ordered as $1 > a > b > a^2 > a^3 > a^4 = 0$. We assume that $a \cdot b = b^2 = a^3$. Then, from this it follows that $a^2 \cdot b = b^3 = 0$. Moreover, (nontrivial) implications are determined as follows. $a \rightarrow b = b \rightarrow a^2 = b \rightarrow a^3 = a$, $\sim b = a^2$, and for $1 \leq m < n \leq 4$, $a^m \rightarrow a^n = a^{n-m}$ if $n - m \neq 2$, and $= b$ otherwise. It is easy to see that K has no nontrivial subalgebras. Thus, it determines an immediate predecessor of Cl which belongs to W_4 .

It will be interesting to see whether there is an immediate predecessor of Cl or not, which is determined by a simple nonlinear FL_{ew} -algebra, or by a simple infinite FL_{ew} -algebra.

In the rest of this section, we will discuss such subdirectly irreducible, *non-simple* FL_{ew} -algebras that determine immediate predecessors of Cl . First, as an application of Theorem 7.2, we will give a series $\{J_n\}_n$ of nonsimple and *nonlinear* FL_{ew} -algebras such that any of them determines an immediate predecessor of Cl . (The same result is obtained independently by T. Kowalski.) For each $n > 1$, the universe of J_n consists of $n + 3$ elements $\{1, a, a^2, \dots, a^n (= a^{n+1}), b, 0\}$. The element b satisfies $a^{n-1} > b > 0$ and is incomparable with

a^n . We assume moreover that $a \cdot b = 0$, and thus $x \cdot b = 0$ for any element $x \leq a$. Then, the following holds:

1. $a^k \rightarrow a^m = a^{m-k}$ if $k < m$,
2. $a^k \rightarrow b = \sim a^k = b$ for $1 \leq k \leq n$,
3. $b \rightarrow a^n = \sim b = a$.

Clearly, the element a generates J_n since $\sim a = b$. Also, b generates it as $\sim b = a$. Finally, each a^k for $1 < k \leq n$ generates it, since $\sim a^k = b$. Thus, J_n has no nontrivial subalgebras. The algebra J_n has the single proper filter $F = \{1, a, a^2, \dots, a^n\}$ except $\{1\}$. It is obvious that the quotient algebra J_n/F is isomorphic to B_2 . Thus, by Theorem 7.2, each J_n determines an immediate predecessor of Cl in W_n . Hence, we have the following.¹³

THEOREM 7.3 *For each $n > 1$, W_n contains an immediate predecessor of Cl , which is determined by the nonsimple and nonlinear FL_{ew} -algebra J_n .*

In the following, we will assume that M is a subdirectly irreducible, non-simple FL_{ew} -algebra satisfying E_n for some $n > 1$, which moreover determines an immediate predecessor of Cl . From Corollary 4.2 it follows that M has the unique coatom a . When $a = 0$, M becomes the 2-valued Boolean algebra, and therefore it is simple, which is a contradiction. Hence, we have that $a > 0$ and $a^n = a^{n+1} > 0$. Let $b = a^n$. Then, the filter F generated by a can be represented as $\{x : b \leq x\}$, which is a proper filter, not equal to $\{1\}$. Therefore, by Theorem 7.2, the quotient algebra M/F must be isomorphic to B_2 . In particular, $\|b\| = \|1\|$ and $\|\sim b\| = \|0\|$ hold, where $\|x\|$ denotes the equivalence class to which an element x belongs. Thus, $b \not\leq \sim b$. Hence, either $\sim b < b$ or b is incomparable with $\sim b$. We assume that the former holds, and consider moreover two special cases; (1) $\sim b = 0$, and (2) $\sim b > 0$ and $\sim \sim b = b$.

1. Case where $\sim b = 0$ holds: Since $b^2 = b$, the set $\{0, b, 1\}$ forms a subalgebra of M , which is isomorphic to the three valued Heyting algebra H_3 . (Note that when M is a Heyting algebra, both $b = a$ and $0 = \sim b < b$ hold always.) Thus, M itself must be isomorphic to H_3 , as it does not contain nontrivial subalgebras.

¹³After the original draft of the present paper is completed, in his master thesis M. Ueda obtained a far stronger result than our Theorem 7.3, which say that for each $n > 1$, W_n contains *countably many* immediate predecessors of Cl , each of which is determined by a nonsimple and nonlinear FL_{ew} -algebra, and that for each $n > 2$, W_n contains also *countably many* immediate predecessors of Cl , each of which is determined by a nonsimple and *linear* FL_{ew} -algebra. Then, T. Kowalski improved them and showed that “countably many” can be replaced by “uncountably many” in them. As Ueda and Kowalski have obtained other stronger results, we have now a much clearer picture of immediate predecessors of Cl and the rest of the present paper should be revised. But, as we mentioned at the beginning, we decided not to make much revisions. For more information on Ueda and Kowalski’s results, see [56] and [40].

2. Case where both $b > \sim b > 0$ and $\sim\sim b = b$ hold: It is easily seen that $(\sim b)^2 = 0$ and $b \rightarrow \sim b = \sim b$ hold. Thus, the set $\{0, \sim b, b, 1\}$ forms a subalgebra of M , which we call l_1 . Hence, M must be isomorphic to l_1 .

From these observations, the following result follows.

THEOREM 7.4 *Suppose that M is a subdirectly irreducible, nonsimple FL_{ew} -algebra satisfying E_n for some $n > 1$ and moreover that $L(M)$ is an immediate predecessor of Cl .*

1. *If M satisfies the axiom $\neg p \vee \neg\neg p$ then M is isomorphic to H_3 .*
2. *If M satisfies both DN and the axiom $(p \rightarrow \neg p) \vee (\neg p \rightarrow p)$ then M is isomorphic to l_1 .*

Proof: Take elements a and b just in the same way as in the above.

- 1). Since $\sim b \vee \sim\sim b = 1$, either $\sim b = 1$ or $\sim\sim b = 1$ holds. If the former is the case, then $a^n = b = 0$. This contradicts our assumption that M is not simple. Thus, $\sim\sim b = 1$ holds and hence $\sim b = 0$ holds. Then, by the above (1) M is isomorphic to H_3 .
- 2). In this case, either $b \leq \sim b$ or $\sim b < b$ holds. But the former does not hold as shown in the above. Since $\sim\sim b = b < 1$ must hold, $\sim b \neq 0$. Thus, M is isomorphic to l_1 by (2). \square

Note that the formula $(\neg p \vee \neg\neg p) \rightarrow ((p \rightarrow \neg p) \vee (\neg p \rightarrow p))$ is provable in FL_{ew} . At present, we do not have any clear view of immediate predecessors of Cl yet, even if we restrict attention only to W_2^{IP} . It is easy to see that $L(L_3)$ is the single member of W_2^{IP} which is determined by a simple FL_{ew} -algebra. The class W_2^{IP} contains $L(l_1)$, $L(j_2)$ and more. In fact, in 2000 T. Kowalski and M. Ueda proved that W_n^{IP} contains uncountably many logics (see [40] for the details).

REFERENCES

- [1] ADILLON, R.J., and V. VERDÚ, On product logic, *Soft Computing* 2, (1998) pp. 141–146.
- [2] AGLIANO, P., I.M.A. FERREIRIM and F. MONTAGNA, Basic hoops: an algebraic study of continuous t-norms, *Studia Logica* 87, (2007) pp. 73–98.
- [3] BAAZ, M., P. HÁJEK, J. KRAJÍČEK and D. ŠVEJDA, Embedding logics into product logic, *Studia Logica* 61, (1998) pp. 35–47.
- [4] BALBES, R. and PH. DWINGER, *Distributive Lattices*, Univ. Missouri Press, 1974.
- [5] BLOK, W.J., The lattice of modal logics: An algebraic investigation, *Journal of Symbolic Logic* 45, (1980) pp. 221–236.

- [6] BLOK, W.J. and I.M.A. FERREIRIM, Hoops and their implicational reducts (abstract), *Algebraic Methods in Logic and Computer Science*, Banach Center Publications 28, (1993) pp. 219–230.
- [7] BLOK, W.J. and I.M.A. FERREIRIM, On the structure of hoops, *Algebra Universalis* 43, (1999) pp. 233–257.
- [8] BLOK, W.J. and D. PIGOZZI, On the structure of varieties with equationally definable principal congruence III, *Algebra Universalis* 32, (1994) pp. 545–608.
- [9] BLOK, W.J. and J.G. RAFTERY, On the quasivariety of BCK-algebras and its subvarieties, *Algebra Universalis* 33, (1995) pp. 68–90.
- [10] BLOK, W.J. and J.G. RAFTERY, Varieties of commutative residuated integral pomonoids and their residuation subreducts, *Journal of Algebra* 190, (1997) pp. 280–328.
- [11] BURRIS, S. and H.P. SANKAPPANAVAR, *A Course in Universal Algebra*, Graduate Texts in Mathematics, Springer, 1981, available on line.
- [12] CASARI, E., Conjoining and disjoining on different levels, *Logic and Scientific Methods*, eds. by M.L. Dalla Chiara et al., Kluwer Academic Publishers, 1997, pp. 261–288.
- [13] CHAGROV, A.V. and M. V. ZAKHARYASCHEV, *Modal Logic*, Oxford University Press, 1997.
- [14] CHANG, C.C., Algebraic analysis of many-valued logics, *Transactions of the American Mathematical Society* 88, (1958) pp. 467–490.
- [15] CHANG, C.C., A new proof of the completeness of the Łukasiewicz axioms, *Transactions of the American Mathematical Society* 93, (1959) pp. 74–80.
- [16] CIGNOLI, R., I.M.L. D’OTTAVIANO and D. MUNDICI, *Algebraic Foundations of Many-valued Reasoning*, Studia Logica Library, Kluwer Academic Publishers, 2000.
- [17] CIGNOLI, R., F. ESTEVA, L. GODO and A. TORRENS, Basic fuzzy logic is the logic of continuous t-norms and their residua, *Soft Computing* 4, (2000) pp. 106–112.
- [18] CORNISH, W.H., Varieties generated by finite BCK-algebras, *Bulletin of Australian Mathematical Society* 22, (1980) pp. 411–430.
- [19] DAVEY, B.A. and H.A. PRIESTLEY, *Introduction to Lattices and Order* 2nd edition, Cambridge University Press, Cambridge, 2002.

- [20] DILWORTH, R.P., Non-commutative residuated lattices, *Transactions of the American Mathematical Society* 46, (1939) pp. 426–444.
- [21] DOŠEN, K. and P. SCHROEDER-HEISTER eds., *Substructural logics*, Oxford Univ. Press, 1993.
- [22] DOŠEN, K., Sequent systems and groupoid models. I, *Studia Logica* 47, (1988) pp. 353–385.
- [23] DOŠEN, K., Sequent systems and groupoid models. II, *Studia Logica* 48, (1989) pp. 41–65.
- [24] FONT, J.M., A.J. RODRÍGUEZ and A. TORRENS, Wajsberg algebras, *Stochastica* 8, (1984) pp. 5–31.
- [25] GALATOS, N., P. JIPSEN, T. KOWALSKI and H. ONO, *Residuated Lattices: an Algebraic Glimpse at Substructural Logics*, Studies in Logic and the Foundations of Mathematics 151, Elsevier, 2007.
- [26] GIRARD, J.Y., Linear logic, *Theoretical Computer Science* 50, (1987) pp. 1–102.
- [27] GRIŠIN, V.N., On a nonstandard logic and its application to the set theory (in Russian), *Issledovanija po formalizovannym jazykam i néklassičéskim logikam*, Moskva, 1974, pp. 135–171.
- [28] GRIŠIN, V.N., On algebraic semantics for logic without the contraction rule (in Russian), *Issledovanija po formalizovannym jazykam i néklassičéskim logikam*, Moskva, 1976, pp. 247–264.
- [29] GRIŠIN, V.N., Predicate and set-theoretical calculi based on logic without the contraction rule (in Russian), *Izvéstia Akademii Nauk SSSR* 45, (1981) pp. 47–68 (also the English translation in *Math. USSR Izvestiya* 18, (1982) pp. 41–59).
- [30] HÁJEK, P., *Metamathematics of Fuzzy Logic*, Trends in Logic 4, Studia Logica Library, Kluwer Academic Publishers, 1998.
- [31] HÁJEK, P., L. GODO and F. ESTEVA, A complete many-valued logic with product-conjunction, *Archive for Mathematical Logic* 35, (1996) pp. 191–208.
- [32] HIGGS, D., Residuated commutative monoids with identity element as least element do not form an equational class, *Mathematica Japonica* 29, (1984) pp. 69–75.
- [33] HÖHLE, U., Commutative residuated monoids, *Non-classical Logics and Their Applications to Fuzzy Subsets* eds. by U. Höhle and P. Klement, Dordrecht, Kluwer, 1995, pp. 53–106.

- [34] IDZIAK, P.M., Lattice operation in BCK-algebras, *Mathematica Japonica* 29, (1984) pp. 839–846.
- [35] ISĒKI, K., BCK-algebras with condition (S), *Mathematica Japonica* 24, (1979) pp. 107–119.
- [36] ISĒKI, K. and S. TANAKA, An introduction to the theory of BCK-algebras, *Mathematica Japonica* 23, (1978) pp. 1–26.
- [37] JANKOV, V.A., On connection between provability in intuitionistic propositional calculus and finite implicative structures (in Russian), *Doklady Akademii Nauk SSSR*, 151, (1963) pp. 1293–1294.
- [38] JOHNSTONE, P.T., Conditions related to de Morgan’s law, *Applications of Sheaves*, Lecture Notes in Mathematics 753, Springer, 1979, pp. 479–491.
- [39] JÓNSSON, B., Algebras whose congruence lattices are distributive, *Mathematica Scandinavica* 21, (1967) pp. 110–121.
- [40] KATOH, Y., T. KOWALSKI and M. UEDA, Almost minimal varieties related to fuzzy logics, *Reports on Mathematical Logic* 41, (2006) pp. 173–194.
- [41] KOMORI, Y., Super-Łukasiewicz propositional logics, *Nagoya Mathematical Journal* 84, (1981) pp. 119–133.
- [42] KOWALSKI, T., The bottom of the lattice of BCK-varieties, *Reports on Mathematical Logic* 29, (1995) pp. 87–93.
- [43] KOWALSKI, T., Semisimplicity, EDPC and discriminator varieties of residuated lattices, *Studia Logica* 77, (2005) pp. 255–265.
- [44] KOWALSKI, T. and H. ONO, The variety of residuated lattices is generated by its finite simple members, *Reports on Mathematical Logic* 34, (2000) pp. 57–75.
- [45] KOWALSKI, T. and H. ONO, Splittings in the variety of residuated lattices, *Algebra Universalis* 44, (2000) pp. 283–298.
- [46] KOWALSKI, T. and H. ONO, *Residuated Lattices: An algebraic glimpse at logics without contraction (preliminary report)*, March, 2001.
- [47] KRULL, W., Axiomatische Begründung der allgemeinen Idealtheorie, *Sitzungsberichte der physikalisch medizinischen Societät der Erlangen* 56, (1924) pp. 47–63.
- [48] MCKENZIE, R.N., G.F. MCNULTY and W.F. TAILOR, *Algebras, Lattices, Varieties* vol.I, Wadsworth & Brooks/Cole, California, 1987.
- [49] MUNDICI, D., MV-algebras are categorically equivalent to bounded commutative BCK-algebras, *Mathematica Japonica* 31, (1986) pp. 889–894.

- [50] OKADA, M. and K. TERUI, The finite model property for various fragments of intuitionistic linear logic, *Journal of Symbolic Logic*, (1999) pp. 790–802.
- [51] ONO, H., Some remarks on semantics for the classical logic without the contraction rules, *Reports on Mathematical Logic* 19, (1985) pp. 3–12.
- [52] ONO, H., Structural rules and a logical hierarchy, *Mathematical Logic*, edited by P.P. PETKOV, Plenum Press, 1990, pp. 95–104.
- [53] ONO, H., Semantics for substructural logics, *Substructural Logics*, eds. by K. Došen and P. Schroeder-Heister, Oxford Univ. Press, 1993, pp. 259–291.
- [54] ONO, H., Proof-theoretic methods in nonclassical logic — an introduction, *Theories of Types and Proofs*, MSJ Memoir 2, eds. by M. Takahashi, M. Okada and M. Dezani-Ciancaglini, Mathematical Society of Japan, 1998, pp. 207–254.
- [55] ONO, H. and Y. KOMORI, Logics without the contraction rule, *Journal of Symbolic Logic* 50 (1985) pp. 169–201.
- [56] ONO, H. and M. UEDA, A classification of logics over FL_{ew} and almost maximal logics, *Philosophical Dimensions of Logic and Science*, eds. by A. Rojczczak, J. Cachro and G. Kurczewski, Kluwer, 2003, pp. 3–13.
- [57] PAŁASIŃSKI, M., Subdirectly irreducibles in E_n -varieties of BCK-algebras, preprint.
- [58] PAVELKA, J., On fuzzy logic II, Enriched residuated lattices and semantics of propositional calculi, *Zeitschrift für der mathematische Logik und Grundlagen der Mathematik* 25, (1979) pp. 119–134.
- [59] ROSENTHAL, K.I., *Quantales and Their Applications*, Pitman Research Notes in Mathematics 234, Longman, 1990.
- [60] SCHWEIZER, B. and A. SKLAR, *Probabilistic Metric Spaces*, North-Holland, 1983.
- [61] SWIRYDOWICZ, K., There exist infinitely many Post-complete extensions of the classical linear logic, manuscript, 1996.
- [62] WARD, M., Residuated distributive lattices, *Duke Mathematical Journal* 6, (1940) pp. 641–651.
- [63] WARD, M. and R.P. DILWORTH, Residuated lattices, *Transactions of the American Mathematical Society* 45, (1939) pp. 335–354.

- [64] WROŃSKI, A., BCK-algebras do not form a variety, *Mathematica Japonica* 28, (1983) pp. 211–213.
- [65] WROŃSKI, A., An algebraic motivation for BCK-algebras, *Mathematica Japonica* 30, (1985) pp. 187–193.
- [66] YUTANI, H., On a system of axioms of commutative BCK-algebras, *Mathematics Seminar Notes of the Kobe University* 5, (1977) pp. 255–256.
- [67] ZLATOŠ, P., Two-levelled logic and model theory, *Colloquia Mathematica Societatis János Bolyai* 28, *Finite Algebra and Multiple-Valued Logic*, Szeged, 1979, pp. 825–872.

The *Australasian Journal of Logic* (ISSN 1448-5052) disseminates articles that significantly advance the study of logic, in its mathematical, philosophical or computational guises. The scope of the journal includes all areas of logic, both pure and applied to topics in philosophy, mathematics, computation, linguistics and the other sciences.

Articles appearing in the journal have been carefully and critically refereed under the responsibility of members of the Editorial Board. Only papers judged to be both significant and excellent are accepted for publication.

The journal is freely available at the journal website at

<http://www.philosophy.unimelb.edu.au/ajl/>.

All issues of the journal are archived electronically at the journal website.

SUBSCRIPTIONS Individuals may subscribe to the journal by sending an email, including a full name, an institutional affiliation and an email address to the managing editor at ajl-editors@unimelb.edu.au. Subscribers will receive email abstracts of accepted papers to an address of their choice. For institutional subscription, please email the managing editor at ajl-editors@unimelb.edu.au.

Complete published papers may be downloaded at the journal's website at <http://www.philosophy.unimelb.edu.au/ajl/>. The journal currently publishes in pdf format.

SUBMISSION The journal accepts submissions of papers electronically. To submit an article for publication, send the \LaTeX source of a submission to a member of the editorial board. For a current list of the editorial board, consult the website.

The copyright of each article remains with the author or authors of that article.